

T.E. (IT)

THEORY OF COMPUTATION

(2015 Course) (Semester - I)

Time : 1 Hour]

[Max. Marks : 30]

Instructions to the candidates:

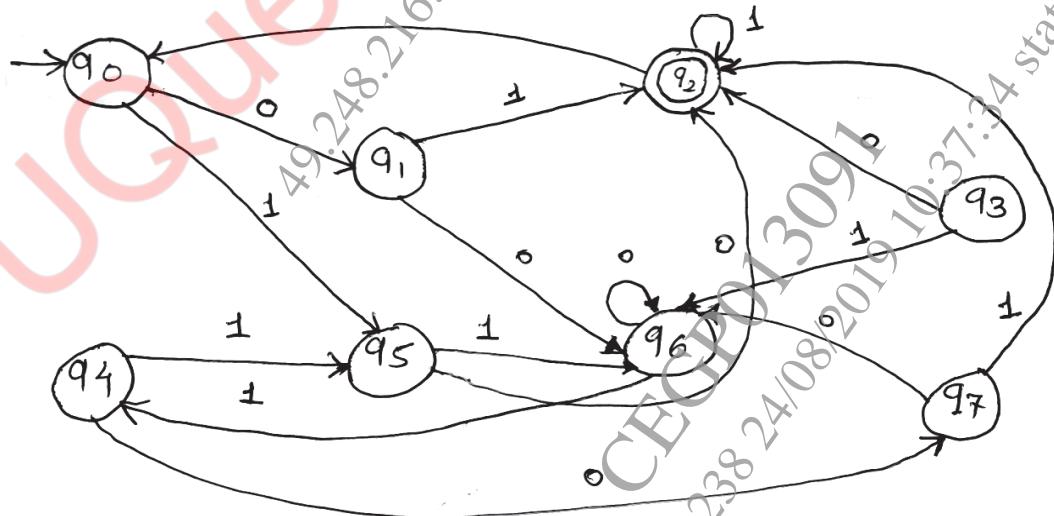
- 1) Attempt questions Q1 or Q2, Q3 or Q4, and Q5 or Q6.
- 2) Neat diagrams must be drawn wherever necessary.
- 3) Figures to the right indicate full marks.
- 4) Assume suitable data, if necessary.

Q1) a) Construct Mealy machine to find 2' complement of any binary number & convert it into Moore machine. [8]

- b)** Define following:-
- i) \in -closure of a state
 - ii) NFA

OR

Q2) a) Construct the minimum state automation equivalent to the transition diagram given below. [8]



- b)** State what do you mean by FSM & state limitations of FSM. [2]

Q3) a) Determine the regular expression over the $\Sigma = \{a, b\}$ for the following [6]

- i) Set of all strings containing exactly 2 a's
- ii) Set of all strings containing at least 2 a's
- iii) Set of all strings that do not consist of two consecutive 0's

b) Convert the given Right-linear Grammar into its equivalent Left-linear Grammar. [4]

$$S \rightarrow 0A|1B$$

$$A \rightarrow 0C|1A|0$$

$$B \rightarrow 1B|1A|1$$

$$C \rightarrow 0|0A$$

OR

Q4) a) Define Pumping Lemma & Apply it to prove the following

$$L = \left\{ 0^{i^2} \mid i \text{ is an integer, } i \geq 1 \right\} \text{ is not regular.} \quad [6]$$

b) Describe in simple english the language defined by the following regular expressions:- [4]

i) $(a+b)^* \cdot aa \cdot (a+b)^*$

ii) $a + b^* \cdot c + \epsilon$

Q5) a) Check whether or not the following Grammar is ambiguous; if it is remove the ambiguity and write an equivalent unambiguous grammar. [6]

$$S \rightarrow i \ C \ t \ S \mid i \ c \ t \ S \ e \ S \mid a$$

$$C \rightarrow b$$

b) Write note on: Chomsky Hierarchy. [4]

OR

Q6) a) Simplify the following Grammar [6]

i) $S \rightarrow aA|bS|\epsilon$

$$A \rightarrow aA|bB|\epsilon$$

$$B \rightarrow aA|bc|\epsilon$$

$$C \rightarrow aC|bc$$

ii) $S \rightarrow A|bb$

$$A \rightarrow B|b$$

$$B \rightarrow S|a$$

b) Find the CNF for the given CFG $S \rightarrow 0S1S|1S0S|\epsilon$ [4]

